

Kan extensions

Topics in category theory – Brno 2018

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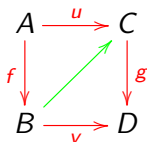
Friday 23rd March, 2018

Many problems in Algebra and Topology can be subsumed by the following general scheme:

$$\begin{array}{ccc} A & \xrightarrow{u} & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{v} & D \end{array}$$

a “calculus of lifting problems and their solution”.

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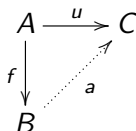
a “calculus of **lifting problems** and their **solution**”.

Motivations

We say that f is left orthogonal to g , or g is right orthogonal to f , if for every such commutative square there is a unique $a : B \rightarrow C$ such that the two triangles commute. This calculus arises, like, everywhere.

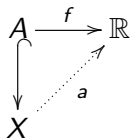
- factorization systems (which you will see) are based on this definition; and
- model categories (if you like algebraic topology look into the definition)

Particularly ubiquitous are problems where $D = *$ is a terminal object, e.g. where we have to find a solution to an extension problem:



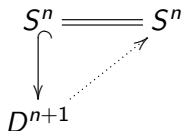
Motivations

- **Tietze extension theorem**: let $f : A \rightarrow \mathbb{R}$ be Lipschitz continuous and defined on a subspace of a metric space $A \subseteq X$; then there exists an extension (non unique) $a : X \rightarrow \mathbb{R}$,



with the same Lipschitz constant of f .

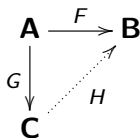
- **Brouwer theorem**, categorically: given a sphere $S^n \subseteq D^{n+1}$



there is **no** continuous retraction $r : D^{n+1} \rightarrow S^n$.

Motivations

To solve a similar problem in the category of categories, i.e. given a diagram

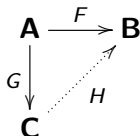


of categories and functors, we face an intrinsic difficulty.

There are many more degrees of freedom given by the fact that in **Cat** there are also **natural transformations** that can fill the diagram **strictly**, **strongly** or **laxly**.

Kan extensions

A reasonable request would be: is there a way to fill



with a natural transformation $\eta : F \Rightarrow HG$? Of course, we do not want to loosen the problem too much, so we might as well ask for some sort of uniqueness to η .

Definition

We can “solve the extension problem at the pair (G, F) ” if there exists a third functor $H : \mathbf{C} \rightarrow \mathbf{B}$ with a natural transformation $\eta : F \rightarrow HG$ which is initial among all such pairs $\langle K, \xi \rangle$.

Kan extensions

This means that

$$\begin{array}{ccc} \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ \downarrow G & & \uparrow \forall \zeta \\ \mathbf{C} & & \end{array} \quad = \quad \begin{array}{ccc} \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ \downarrow G & \nearrow \eta & \uparrow \exists \bar{\zeta} \\ \mathbf{C} & \nearrow H & \end{array} \quad \begin{array}{c} \\ \\ \end{array} \quad \begin{array}{c} \\ \\ K \end{array}$$

Definition

We say that the pair $\langle H, \eta \rangle$ “**exhibits** $\text{Lan}_G F$ ” when it has this property; we often say simply that “ H is $\text{Lan}_G F$ ”.

pointwise Kan extensions

It turns out that this definition, albeit correct, is **too general**.

To save the situation we can give a less general one encompassing many (actually, all) examples of interest. This leads to the definition of **pointwise Kan extension**

Idea

It would be nice to compute $\text{Lan}_G F$ “locally”, in the following way: given $c \in \mathbf{C}$

$$\begin{array}{ccc} \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ \downarrow G & \Downarrow & \nearrow \text{Lan}_G F \\ * & \xrightarrow{[c]} & \mathbf{C} \end{array}$$

pointwise Kan extensions

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$$\begin{array}{ccccc} (G \downarrow c) & \xrightarrow{\Sigma} & \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ & \searrow \Downarrow & \downarrow G & \Downarrow & \nearrow \text{Lan}_G F \\ & & * & \xrightarrow{[c]} & \mathbf{C} \end{array}$$

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$$\begin{array}{ccccc} (G \downarrow c) & \xrightarrow{\Sigma} & \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ & \searrow \swarrow & \downarrow G & \swarrow \nwarrow & \\ & & \mathbf{C} & \xrightarrow{\text{Lan}_G F} & \mathbf{B} \\ & \downarrow & \uparrow [c] & & \\ * & \xrightarrow{\quad} & \mathbf{C} & & \end{array}$$

and $\text{Lan}_G F(c) = \text{Lan}_G F \circ [c] \cong \text{colim}_{(G \downarrow c)} F \circ \Sigma$.

It turns out that this simple idea actually works.

Proposition

When the colimit

$$\operatorname{colim}_{(G \downarrow c)} F \circ \Sigma$$

exists for every $c \in \mathbf{C}$, the assignment $c \mapsto \operatorname{colim}_{(G \downarrow c)} F \circ \Sigma$ can be promoted to a functor $\mathbf{C} \rightarrow \mathbf{B}$ that has the same universal property of $\operatorname{Lan}_G F$ in the category $[\mathbf{C}, \mathbf{B}]$.

This special instance of $\operatorname{Lan}_G F$ is called the “pointwise” left Kan extension of F along G , to reinforce the intuition that the value of $\operatorname{Lan}_G F$ can be computed “on points” of its domain.

Definition

A functor $P : \mathbf{B} \rightarrow \mathbf{X}$ preserves $\text{Lan}_G F$ if $P \circ \text{Lan}_G F \cong \text{Lan}_G PF$.

Proposition

Let $\mathbf{X} = \mathbf{Set}$ and $P = \mathbf{B}(-, X) : \mathbf{B} \rightarrow \mathbf{Set}$ a representable presheaf. Then P preserves every pointwise Kan extension.

$$\begin{aligned} P \circ \text{Lan}_G F &\cong \text{hom}(\text{Lan}_G F(C), X) \cong \text{hom}(\text{colim}_{A:G \downarrow c} FA, X) \\ &\cong \lim_{A:G \downarrow c} \text{hom}(FA, X) \cong \lim_{A:G \downarrow c} P(FA) \cong \text{Lan}_G PFC \end{aligned}$$

Is the converse true?

A fundamental theorem

Let $\text{Lan}_G F$ be pointwise. The the correspondence $F \mapsto \text{Lan}_G F$ can be promoted to a functor

$$\text{Lan}_G : [\mathbf{A}, \mathbf{B}] \rightarrow [\mathbf{C}, \mathbf{B}]$$

and there is an adjunction $\text{Lan}_G \dashv G^*$ (where $G^*(K) = K \circ G$).

There are many ways to prove this statement; we opt for an hands-on proof, defining a natural bijection

$$[\mathbf{C}, \mathbf{B}](\text{Lan}_G F, H) \leftrightarrow [\mathbf{A}, \mathbf{B}](F, HG)$$

proving also that η is the unit of the adjunction.

Haskell-y remark

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Lan_G    :: (A -> B) -> (C -> B)  --this is a left adjoint
Lan_G F  :: C -> B  --this is not (always)
```

A fundamental theorem

Exercise

Prove the above proposition.

Exercise

Prove the dual of the above “fundamental theorem” for **right** Kan extensions, defined below.

Exercise

What is the colimit expressing the pointwise extension $\text{Lan}_1 F$? Relate this result with the Yoneda lemma.

Right Kan extensions

- A **right** Kan extension of F along G is a pair $\langle \text{Ran}_G F : \mathbf{C} \rightarrow \mathbf{B}, \epsilon : \text{Ran}_G F \circ G \rightarrow F \rangle$ terminal among all such pairs.
- A pointwise right Kan extension is obtained from the limit

$$\begin{array}{ccccc}
 (c \downarrow G) & \xrightarrow{\Sigma} & \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\
 \downarrow & \nearrow & \downarrow G & \nearrow & \\
 * & \xrightarrow{[c]} & \mathbf{C} & & \\
 & & & \nearrow \text{Ran}_G F &
 \end{array}$$

- There is an adjunction $G^* \dashv \text{Ran}_G$ with counit ϵ .

In nice situations (for example when \mathbf{B} is complete and cocomplete) every functor $G : \mathbf{A} \rightarrow \mathbf{C}$ defines a triple of adjoints

$$\text{Lan}_G \dashv G^* \dashv \text{Ran}_G : [\mathbf{C}, \mathbf{B}] \begin{array}{c} \xleftarrow{\text{Ran}_G} \\ \xrightarrow{\quad} \\ \xleftarrow{\text{Lan}_G} \end{array} [\mathbf{A}, \mathbf{B}]$$

Examples

Limits and colimits

All colimits are Kan extensions: let $\mathbf{C} = *$ be the terminal category and $G = ! : \mathbf{A} \rightarrow *$ the terminal functor.

Then the functor $(G \downarrow c) \rightarrow \mathbf{A}$ is an isomorphism and

$$\operatorname{colim}_{(! \downarrow c)} F \circ \Sigma \cong \operatorname{colim} F$$

(because Σ is an equivalence; exercise: prove that an equivalence is a *cofinal* functor).

$$\begin{array}{ccccc} (G \downarrow c) & \xrightarrow{\Sigma} & \mathbf{A} & \xrightarrow{F} & \mathbf{B} \\ \downarrow & \Downarrow & \downarrow G & \Downarrow & \nearrow \operatorname{Lan}_G F \\ * & \xrightarrow{[c]} & \mathbf{C} & & \end{array}$$

Dually, $\lim_{(c \downarrow !)} F \Sigma \cong \lim F$.

Adjoints

Under reasonable conditions, Kan extensions provide adjoints to functors.

Proposition

Let $R : \mathbf{C} \rightarrow \mathbf{D}$ be a functor. The following conditions are equivalent

- R has a left adjoint L with unit and counit $L \frac{\eta}{\epsilon} \mid R$.
- $\langle L, \epsilon \rangle$ exhibits $\text{Ran}_R 1$ and R preserves this right Kan extension

The proof can be done elementarily (good luck) or with a moderate amount of 2-cell calculus, recalling the zigzag identities

$$\begin{array}{ccc} \mathbf{C} & \xlongequal{\quad} & \mathbf{C} \\ R \searrow & & \nearrow L \\ & \mathbf{D} & \\ & \xlongequal{\quad} & \mathbf{D} \\ & & \searrow R \end{array} = 1_R \quad ; \quad \begin{array}{ccc} & \mathbf{C} & \\ L \nearrow & \xlongequal{\quad} & \searrow R \\ \mathbf{D} & \xlongequal{\quad} & \mathbf{D} \\ & & \nearrow L \end{array} = 1_L$$

for the unit and counit of an adjunction $L \frac{\eta}{\epsilon} \mid R$.

Let $i : H \leq G$ be the inclusion of a subgroup regarded as a functor between one-object categories; if k is a field, the category of k -linear representations of G can be identified with the functor category $[G, \mathbf{Vect}_k]$, and similarly for H .

The category of finite dimensional vector spaces has finite co/limits, so if G, H are finite the functor

$$[H, \mathbf{Vect}] \rightarrow [G, \mathbf{Vect}]$$

induced by i has both a left and a right adjoint by virtue of our fundamental theorem.

Induction and coinduction

These adjoints, Lan_i and Ran_i , can be respectively computed as (V is a H -module)

$$V \mapsto k[G] \otimes_{k[H]} V \in G\text{-Vect} \quad V \mapsto \text{hom}_{k[H]}(k[G], V) \in G\text{-Vect}$$

These, in turn, can be computed as the limit and colimit, resp.,

$$\begin{aligned} \text{hom}_{k[H]}(k[G], V) &\longrightarrow \text{hom}_k(K[G], V) \begin{array}{c} \xrightarrow{\quad} \\ \xrightarrow{\quad} \end{array} \prod_{h \in H} \text{hom}(k[G], V) \\ \coprod_{h \in H} K[G] \otimes_k V &\begin{array}{c} \xrightarrow{\quad} \\ \xrightarrow{\quad} \end{array} k[G] \otimes_k V \longrightarrow k[G] \otimes_{k[H]} V \end{aligned}$$

Where we take the equivariant linear maps of $k[H]$ -modules and we quotient out for the relation prescribing $(h.\alpha) \otimes v - \alpha \otimes (h.v)$.

Let X be a topological space and $Op(X)$ the category of its open subsets. A presheaf on X is a functor $Op(X)^{op} \rightarrow \mathbf{Set}$. The category of sets is complete and cocomplete, so that every continuous function $f : Y \rightarrow X$ induces a functor

$$[Op(X)^{op}, \mathbf{Set}] \xrightarrow{f^*} [Op(Y)^{op}, \mathbf{Set}]$$

that has both a left and a right adjoint: $f_! \dashv f^* \dashv f_*$.

These adjoints are computed as follows:

$$\begin{cases} f_!(\mathcal{F})(V) = \operatorname{colim}_{U \supseteq fV} \mathcal{F}(U) \\ f_*(\mathcal{F})(V) = \mathcal{F}(f^{\leftarrow} V) \end{cases}$$

Exercise: the first equation is the formula of our fundamental theorem: prove it.

A **directed graph** is a covariant functor on the category $\Gamma = \{E \begin{smallmatrix} \xrightarrow{s} \\ \xrightarrow{t} \end{smallmatrix} V\}$ (a directed graph is a set of vertices, a set of edges, and two functions giving source and target). The “forgetful” functor $U : \mathbf{dGph} \rightarrow \mathbf{Set}$ is the functor

$$t_{\Gamma}^* : [\Gamma, \mathbf{Set}] \rightarrow \mathbf{Set} = [* , \mathbf{Set}]$$

where the functor $t : * \rightarrow \Gamma$ chooses the object V and its identity morphism. Since \mathbf{Set} is co/complete, t_{Γ}^* has both a left and a right adjoint by virtue of our fundamental theorem.

Exercise

Show that $\text{Lan}_{t_{\Gamma}}$ exists and it is given by the functor that sends a set X to the directed graph having X as set of vertices, and no edges.

What does the functor $\text{Ran}_{t_{\Gamma}}$ do?

Let $\mathbf{\Delta}$ be the category of nonempty finite totally ordered sets and monotone functions; denote $[n] = \{0 < 1 < \dots < n\}$ the generic object and $\Delta[n]$ the representable object on it. Let $J : \mathbf{\Delta} \rightarrow \mathbf{Cat}$ the tautological functors that regards a poset as a category.

We can define a functor $\mathbf{Cat} \rightarrow [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}]$ (the category of **simplicial sets**) such that $[n]$ is sent to

$$N(\mathcal{C})_n = \mathbf{Cat}(J[n], \mathcal{C}) = \{C_0 \rightarrow C_1 \rightarrow \dots \rightarrow C_n \mid C_i \in \mathcal{C}\}$$

Proposition

The functor $\mathcal{C} \mapsto (\lambda[n].N(\mathcal{C})_n)$ has a left adjoint given by $\text{Lan}_y J : [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}] \rightarrow \mathbf{Cat}$, where $y : \mathbf{\Delta} \rightarrow [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}]$ is the Yoneda embedding functor.

We prove that $\text{Lan}_y J \dashv N$ by providing a natural bijection of hom-sets: given a category \mathcal{D} and a simplicial set X we have

$$\begin{aligned}
 \mathbf{Cat}(\text{Lan}_y J(X), \mathcal{D}) &\cong \mathbf{Cat} \left(\text{colim}_{([n], x) \in (y \downarrow X)} J\Sigma([n], x), \mathcal{D} \right) \\
 &\cong \lim_{([n], x) \in (y \downarrow X)} \mathbf{Cat}(J\Sigma([n], x), \mathcal{D}) \\
 &\cong \lim_{([n], x) \in (y \downarrow X)} \mathbf{Cat}(J[n], \mathcal{D}) \\
 &\cong \lim_{([n], x) \in (y \downarrow X)} [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}](\Delta[n], N(\mathcal{D})) \\
 &\cong [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}](\text{colim}_{([n], x) \in (y \downarrow X)} \Delta[n], N(\mathcal{D})) \\
 &\cong [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}](X, N(\mathcal{D}))
 \end{aligned}$$

By virtue of a similar argument it is possible to prove that there is an adjunction

$$\begin{array}{ccc} \mathbf{\Delta} & \xrightarrow{\rho} & \mathbf{Top} \\ y \downarrow & \nearrow \text{Lan}_y \rho & \\ [\mathbf{\Delta}^{\text{op}}, \mathbf{Set}] & \xleftarrow{S} & \end{array}$$

where the functor $S(X)$ is the simplicial set $[n] \mapsto \mathbf{Top}(\rho[n], X)$ and $\rho[n]$ is the topological simplex

$$\rho[n] = \{\vec{X} \in \mathbb{R}^n \mid X_i \geq 0, \sum X_i \leq 1\}$$

Of course there is a pattern here:

Theorem

Let

$$\begin{array}{ccc} \mathbf{C} & \xrightarrow{F} & \mathbf{D} \\ y \downarrow & & \\ [\mathbf{C}^{\text{op}}, \mathbf{Set}] & & \end{array}$$

be an arrangement of functors where \mathbf{C} is small and \mathbf{D} is cocomplete.

Of course there is a pattern here:

Theorem

Let

$$\begin{array}{ccc} \mathbf{C} & \xrightarrow{F} & \mathbf{D} \\ y \downarrow & \text{Lan}_y F \nearrow & \\ [\mathbf{C}^{\text{op}}, \mathbf{Set}] & & N_F \searrow \end{array}$$

be an arrangement of functors where \mathbf{C} is small and \mathbf{D} is cocomplete. Then there is an adjunction $\text{Lan}_y F \dashv N_F = \text{hom}(F, 1)$ where F commutes with finite limits if and only if $\text{Lan}_y F$ does.

Examples of nerves: Isbell duality

Isbell duality sets up an adjunction

$$[\mathbf{A}^{\text{op}}, \mathbf{Set}] \underset{\text{Spec}}{\overset{\mathcal{O}}{\rightleftarrows}} [\mathbf{A}, \mathbf{Set}]^{\text{op}}$$

for every small category \mathbf{A} , defined as

$$\mathcal{O} : X \mapsto [\mathbf{A}^{\text{op}}, \mathbf{Set}](X, \text{hom}(c, -))$$

$$\text{Spec} : Y \mapsto [\mathbf{A}, \mathbf{Set}](Y, \text{hom}(-, c))$$

Riddle

Why are the functors called \mathcal{O} (“ring of functions”) and Spec (“spectrum”)?

Examples of nerves: Isbell duality

Exercise

Show that

$$\begin{array}{ccc} & \mathbf{A} & \\ y \swarrow & & \searrow y^\circ \\ [\mathbf{A}^{\text{op}}, \mathbf{Set}] & \begin{array}{c} \xrightarrow{\text{Lan}_y(y^\circ)} \\ \perp \\ \xleftarrow{\text{Lan}_{y^\circ}(y)} \end{array} & [\mathbf{A}, \mathbf{Set}]^{\text{op}} \end{array}$$

is the Isbell adjunction.

Examples of nerves: space from sheaves

Let X be a topological space and $Op(X)$ the category of its open subsets. Define the tautological functor

$$\begin{aligned} Op(X) &\rightarrow \mathbf{Top}/X \\ U &\mapsto (U \subseteq X) \end{aligned}$$

Since \mathbf{Top}/X is cocomplete, we can define an adjunction

$$\begin{array}{ccc} [Op(X)^{op}, \mathbf{Set}] & \begin{array}{c} \xrightarrow{\text{etale}} \\ \xleftarrow{\text{sheafofsections}} \end{array} & \mathbf{Top}/X \end{array}$$

$$\begin{array}{ccc} U & \overset{\dots\dots\dots}{\longrightarrow} & E \\ \searrow i & & \swarrow p \\ & X & \end{array} \quad \leftarrow \quad \left[\begin{array}{c} E \\ \downarrow \\ X \end{array} \right]$$

that suitably restricted and corestricted determines the equivalence between *sheaves* and étale spaces over X .

A few references

- 3.7 of F. Borceux “[Handbook of categorical algebra](#)”.
- Kashiwara and Schapira, section 2.3 in “[Categories and Sheaves](#)” (for an hands-on proof of the fundamental theorem).
- M. Vergura’s [notes](#) (link) on Kan extensions.
- M. Lehner’s B.Sc. [thesis](#) (link) “[All concepts are Kan extensions](#)”
- I.1–5 of Dubuc’s “[Kan extensions in enriched category theory](#)”.