

The art of \int – notable integrals in Category Theory

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ItaCa Liber I

Plan de l'œuvre

«I have always disliked analysis»

P.J. Freyd (Algebraic real analysis)

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Caveat: this wants to be a "light" talk (and partly self-promotion).

Coends

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Coends are universal objects associated to functors

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$$\int_{\mathcal{C}} T \xrightarrow{\text{terminal}} T(X, X) \quad T(X, X) \xrightarrow{\text{initial}} \int^{\mathcal{C}} T$$

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$$A \otimes_G B \cong \text{colim} \left(\bigoplus_{g \in G} A \otimes B \begin{array}{c} \xrightarrow{g \otimes 1} \\ \xrightarrow{1 \otimes g} \end{array} A \otimes B \right)$$

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$$\text{hom}_G(A, B) \cong \lim \left(\text{hom}(A, B) \rightrightarrows \prod_{g \in G} \text{hom}(A, B) \right)$$

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- They satisfy a **Fubini rule**

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- They provide analogues for a **theory of integrations**

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Analysis

Dirac deltas

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Every object of the form yC is **tiny**, so it is a functor concentrated on the “point” $C \in \mathcal{C}$. **Yoneda lemma** says that

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Or dually,

$$\int^C yC(X) \times FX = yC \otimes_{\mathcal{C}} F \cong FC$$

(the **Dirac δ functor** concentrated on $C \in \mathcal{C}$ evaluates functors on points).

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the exterior derivative.

- Let $C : \mathbf{N}^{\text{op}} \rightarrow \text{Mod}(\mathbf{R})$ the functor sending n to (the vector space over) smooth maps $Y \rightarrow X$ where Y is closed n -dimensional oriented manifold: $\partial : C_{n+1} \rightarrow C_n$ is the geometric boundary.

Stokes' theorem

$$C \otimes \Omega : \mathbf{N}^{\text{op}} \times \mathbf{N} \rightarrow \mathbf{Vect}$$

is a functor. There is a map

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$$(Y \xrightarrow{\varphi} X, \omega) \mapsto \int_Y \varphi^* \omega$$

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Theorem (Stokes): The square

$$\begin{array}{ccc} C_{n+1} \otimes \Omega_n & \xrightarrow{\partial \otimes 1} & C_n \otimes \Omega_n \\ \downarrow 1 \otimes d & & \downarrow \int_n \\ C_{n+1} \otimes \Omega_{n+1} & \xrightarrow{\int_{n+1}} & \mathbf{R} \end{array}$$

is commutative for every $n \in \mathbf{N}$. $\int^n C \otimes \Omega$ is a certain $H^0 \dots$

Distributions

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A **Lawvere distribution** is a left adjoint between two toposes;

dist. between sheaves on \mathcal{C}, \mathcal{D}
||
profunctors $p : \mathcal{C} \rightsquigarrow \mathcal{D}$

(Dirac distributions over a topos \mathcal{E} are **points** of that topos (geometric morphisms $p : \mathcal{E} \rightarrow \mathbf{Set}$); complies with intuition)

Convolutions

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$$f * g = y \mapsto \int_G f(x) \cdot g(y - x) d\mu_G$$

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Let \mathcal{C} be a monoidal category; the category $[\mathcal{C}, \mathbf{Set}]$ becomes a monoidal category with a convolution operation of presheaves:

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(what if \mathcal{C} is closed? You recover the above formula)

Fourier theory

Fact:

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A profunctor $K : \mathcal{C} \rightsquigarrow \mathcal{D}$ between monoidal categories is a **multiplicative kernel** if the associated

$$\hat{K} : [\mathcal{C}^{\text{op}}, \mathcal{V}] \rightleftarrows [\mathcal{D}^{\text{op}}, \mathcal{V}]$$

is a strong monoidal adjunction wrt convolution product.

Fourier theory

The **K -Fourier transform** $f \mapsto \mathfrak{F}_K(f) : \mathcal{D} \rightarrow \mathcal{V}$, obtained as the image of $f : \mathcal{C} \rightarrow \mathcal{V}$ under the left Kan extension $\text{Lan}_y K : [\mathcal{C}, \mathcal{V}] \rightarrow [\mathcal{D}, \mathcal{V}]$.

$$\mathfrak{F}_K(f) : X \mapsto \int^A K(A, X) \otimes fA.$$

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The **dual Fourier transform** is defined as:

$$\mathfrak{F}^\vee(g) : Y \mapsto \int_A [K(A, X), gA]$$

(prove the relation $\mathfrak{F}_K^\vee(g) \cong \mathfrak{F}_K(g^*)^*$)

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\mathfrak{F}_y is the identity functor; analogue in analysis, what is the Fourier transform of δ ?

Fourier theory

- \mathfrak{F}_K preserves the **upper convolution** of presheaves f, g , defined as

$$f \bar{*} g = \int^{AA'} fA \otimes gA' \otimes \mathcal{C}(A \otimes A', -);$$

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dually,

- \mathfrak{F}_K^\vee preserves the **lower convolution** of presheaves f, g , defined as

$$f \underline{*} g = \int_{AA'} (fA^* \otimes (gA')^* \otimes \mathcal{C}(A \otimes A', -))^*$$

Fourier theory

Define the **pairing** $(\mathcal{C}, \mathcal{V}) \times (\mathcal{C}, \mathcal{V}) \rightarrow \mathcal{V}$ as the twisted form of functor tensor product

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If K is a kernel such that $\text{Lan}_y K$ is fully faithful, we have **Parseval formula**:

$$\langle f, g \rangle \cong \langle \mathfrak{F}_K(f), \mathfrak{F}_K(g) \rangle.$$

Bibliography

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- Yoneda, Nobuo. **On Ext and exact sequences**. J. Fac. Sci. Univ. Tokyo Sect. I 8.507-576 (1960): 1960.

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- Day, Brian J. **Monoidal functor categories and graphic Fourier transforms**. Theory and Applications of Categories 25.5 (2011): 118-141.

When you come across a paper with page after page of nothing but enriched categories and coend formulas:

