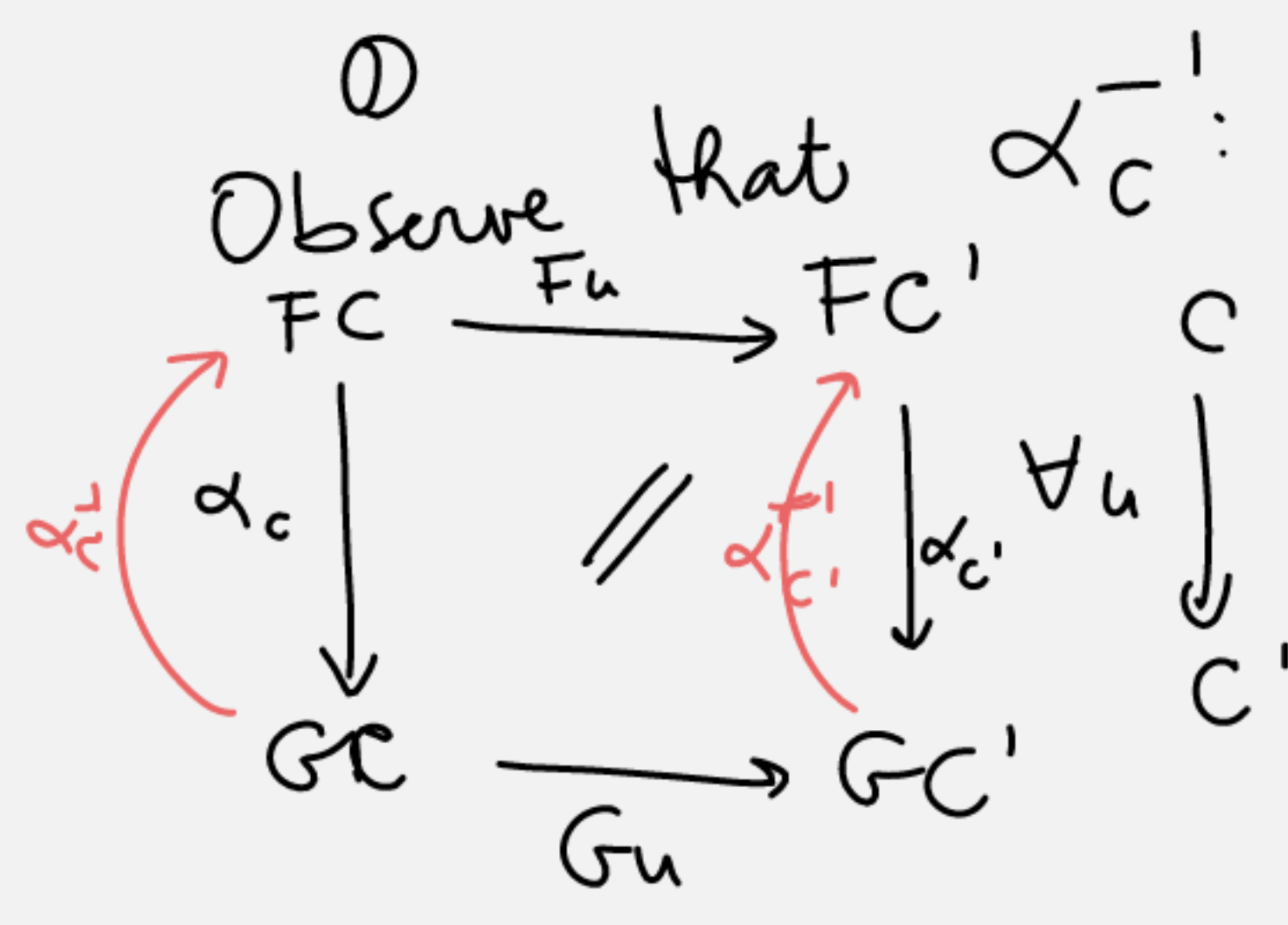


$\alpha: F \Rightarrow G$ natural
 is a natural isomorphism iff all its components are iso (in the codomain of F, G)

$\mathcal{D} \left(\begin{array}{c} \mathcal{C} \\ \xrightarrow{\alpha} \\ \mathcal{C} \end{array} \right) \mathcal{D}$ $\alpha_c: FC \longrightarrow GC$ are morphisms in \mathcal{D}

$\alpha^{-1}: G \Rightarrow F$ is natural

(inverses of α_c)



$$\alpha_{c'} \circ F_u = G_u \circ \alpha_c$$

$$\alpha_{c'}^{-1} \circ \alpha_{c'} \circ F_u \circ \alpha_c^{-1} = \alpha_{c'}^{-1} \circ G_u \circ \underbrace{\alpha_c \circ \alpha_c^{-1}}_1$$

$$\underbrace{1}_{\alpha_c^{-1} \circ \alpha_c} \circ F_u \circ \alpha_c^{-1} = \alpha_{c'}^{-1} \circ G_u \circ 1$$

$$F_u \circ \alpha_c^{-1} = \alpha_{c'}^{-1} \circ G_u \quad \text{is the naturality condition for } \alpha^{-1}$$

Definition $\mathcal{C} \xrightarrow{F} \text{Set}$ is called a PRESHEAF

Given any object of \mathcal{C} , $X \in \mathcal{C}_0$
 I can consider the functor $\text{hom}(-, X)$ def'd as follows

for $A \in \mathcal{C}_0$ $\text{hom}(-, X)(A) = \text{hom}(A, X) \in \text{Set}$
 $\{A \xrightarrow{f} X\}$

$u: A \rightarrow B$ $\text{hom}(u, X): \text{hom}(A, X) \leftarrow \text{hom}(B, X)$
 $\xrightarrow{- \circ u} \left(A \xrightarrow{f \circ u} X \right) \leftarrow f (f: B \rightarrow X)$

This is a functor 'bc. \mathcal{C} is a category":

$\text{hom}(id_A, X) = id_{\text{hom}(A, X)}$ $f \circ id = f$ by category axiom

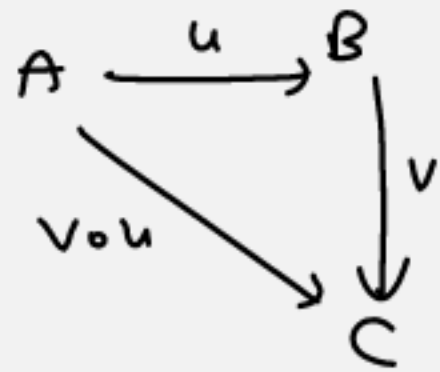
$\text{hom}(v \circ u, X) = \text{hom}(u, X) \circ \text{hom}(v, X)$

$\text{hom}(v \circ u, X)(f) = f \circ (v \circ u)$

$\text{hom}(u, X)(\text{hom}(v, X)(f))$

$\text{hom}(u, X)(f \circ v)$

$(f \circ v) \circ u$



Def a functor $F: \mathcal{C}^{op} \rightarrow \text{Set}$ with the property that there exists an invertible natural transformation

$$\theta: \text{hom}(-, X_F) \xrightarrow{\cong} F$$

REPRESENTABLE
functor.

for some object $X_F \in \mathcal{C}_0$, is called a

(It is a non-obvious fact that when X_F exists, it is unique.)
it's a consequence of "Yoneda lemma")

$$\text{Cat} \longrightarrow \text{Set}$$

$$\mathcal{C} \longmapsto \mathcal{C}_0$$

$$\mathcal{C} \longmapsto \mathcal{C}_1$$

$$\mathcal{C} \longmapsto \mathcal{C}_2 = \{ X \xrightarrow{f} Y \xrightarrow{g} Z \}$$

$$\mathcal{C} \longmapsto \mathcal{C}_3 = \{ \longrightarrow \longrightarrow \longrightarrow \}$$

⋮

$$\mathcal{C} \longmapsto \mathcal{C}_n = \{ \longrightarrow \longrightarrow \longrightarrow \dots \longrightarrow \}$$

$$\text{Top}^{op} \longrightarrow \text{Set}$$

$$(X, \tau) \longmapsto \tau \subseteq 2^X$$

∃ topological space S
w/ the property

$$\left\{ \begin{array}{l} \text{continuous function} \\ X \rightarrow S \end{array} \right\} \xleftrightarrow{\cong} \left\{ \begin{array}{l} \text{open} \\ \text{subsets} \\ \text{of } X \end{array} \right\}$$

$$S = \{ a, b \}$$

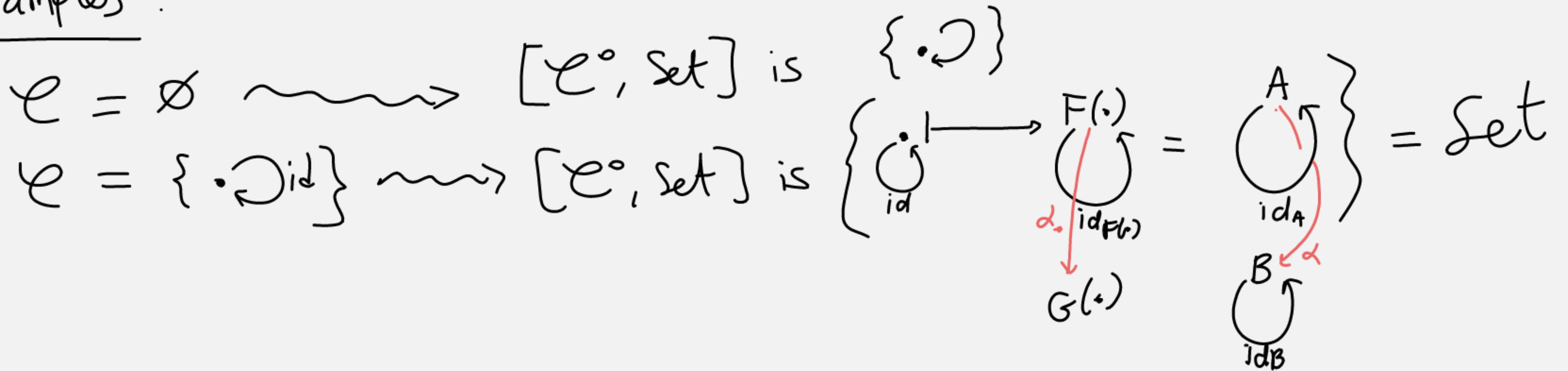
$$\{ \emptyset, \{a\}, \{a, b\} \}$$

There exists a category $[\mathcal{C}^{\circ}, \text{Set}]$
of presheaves (= contravariant functors from \mathcal{C} to Set)
and natural transformations

(This is what made the previous def'n sensible)

bc natural isos are exactly isomorphisms $[\mathcal{C}^{\circ}, \text{Set}]$

Examples:



$$\mathcal{C} = \{0, 1\}$$

$\mathcal{C}^{\text{op}} \xrightarrow{F} \text{Set}$ sends $0 \mapsto$ a set $F(0)$

$1 \mapsto$ a set $F(1)$

(2 identities to identities)

$[\mathcal{C}^{\text{op}}, \text{Set}]$ is the category $\text{Set} \times \text{Set}$

$F \mapsto (F_0, F_1)$

$F_{AB} \mid \begin{array}{l} 0 \mapsto A \\ 1 \mapsto B \end{array} \leftarrow (A, B)$

Rmk Observe that this argument works more generally

if $\mathcal{C} = \text{discrete on a set } A$, $\{\text{functions } \mathcal{C} \xrightarrow{F} \mathcal{D}\} \equiv$

More generally regarding A^{set} as a discrete category

$$[A^{\text{op}}, \text{Set}] \equiv \text{Set}^A$$



$\{\text{functions } A \xrightarrow{F_0} \mathcal{D}_0\}$
objects of \mathcal{D}

$$\mathcal{C} = \{0 \xrightarrow{s} 1\}$$

$$\mathcal{C}^{op} \xrightarrow{G} \text{Set}$$

$$G(0) = \text{Vertices}$$

directed graph

$$d\text{Gph} = [\{0 \rightrightarrows 1\}^{op}, \text{Set}]$$

$$G_s \uparrow \uparrow G_t$$

$$G(1) = \text{Edges}$$

$$\text{hom}(-, 0) : \{0 \rightrightarrows 1\}^{op} \rightarrow \text{Set}$$

$$\text{hom}(-, 1) : \quad \quad \quad "$$

What is a representable diagram?

$$\begin{aligned} \text{hom}(-, 0)(0) &= \text{hom}(0, 0) = \{id_0\} \\ (1) &= \text{hom}(1, 0) = \{\cancel{\quad}\} \emptyset \end{aligned}$$

Γ_0

$$\{ \bullet \xrightarrow{id_0} \bullet \}$$

$$\begin{aligned} \text{hom}(-, 1)(0) &= \text{hom}(0, 1) = \{s, t\} \\ \text{hom}(-, 1)(1) &= \text{hom}(1, 1) = \{id_1\} \end{aligned}$$

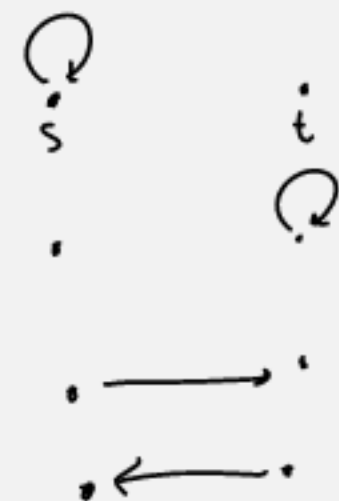
Γ_1

$$\text{hom}(s, 1) = \sigma \uparrow \uparrow \tau = \text{hom}(t, 1)$$

$$= \Gamma_1(1) = \{ \underset{x}{\bullet} \xrightarrow{id_1} \bullet \}$$

$$\sigma(x) = x \circ s = id \circ s = s$$

$$\tau(x) = x \circ t = id \circ t = t$$



(Yoneda lemma: $G \in d\text{Gph}$)

$$G \iff \begin{aligned} \alpha_0: \Gamma_0 \Rightarrow G \\ \alpha_1: \Gamma_1 \Rightarrow G \end{aligned} \quad G: \{0 \rightrightarrows 1\}^{op} \rightarrow \text{Set}$$

$$\alpha_0: \{ \bullet \} \rightarrow G$$

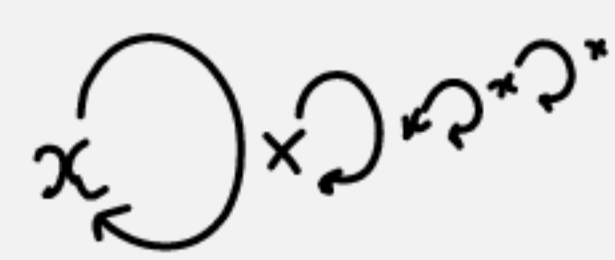
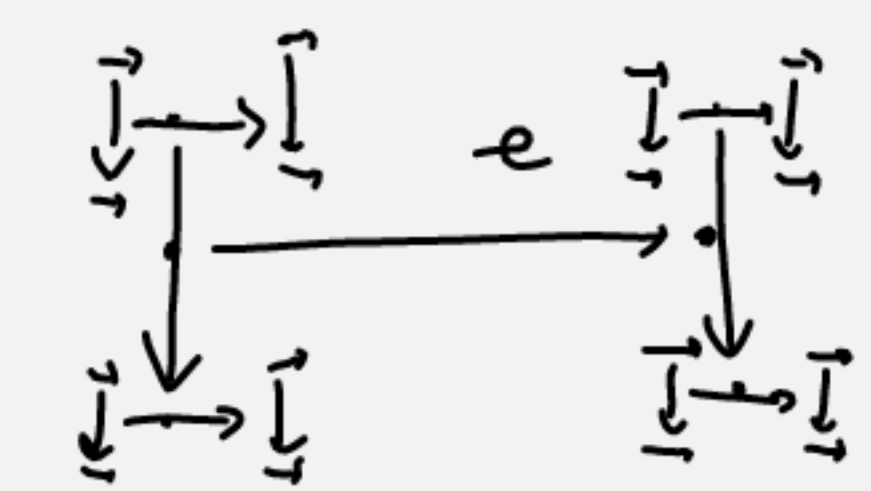
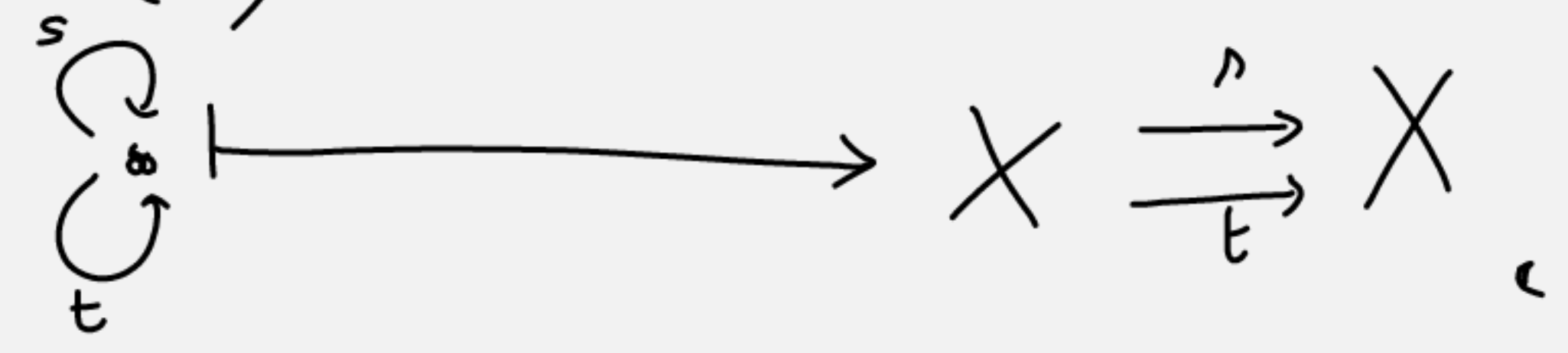
$$\alpha_1: \{ \bullet \rightarrow \bullet \} \rightarrow G$$



$F(2) = F(s, t)$
free monoid
on $\{s, t\}$

$$= \{ [], [s], [t], [st], [ts], [sst], [sts], [stt], \dots \}$$

$X: F(2) \longrightarrow \text{Set}$



$X_0 \subset X$
 $X_0 = \{x \in X \mid s(x) = x = t(x)\}$