

$$\begin{array}{l}
 \mathcal{C} \xrightarrow{\mathcal{A}} [\mathcal{C}^0, \text{Set}] \\
 \mathcal{C}^0 \xrightarrow{\text{co}\mathcal{A}} [\mathcal{C}, \text{Set}] \\
 \mathcal{C}^0 \times \mathcal{C} \xrightarrow{\text{hom}} \text{Set}
 \end{array}
 \quad
 \begin{array}{l}
 X \mapsto \mathcal{C}(-, X) (= \text{hom}(-, X)) \\
 A \mapsto \mathcal{C}(A, -) (= \text{hom}(A, -))
 \end{array}$$

hom functor

Yoneda & coYoneda arise from "saturating" hom in one or the other variable

$$\begin{array}{ccc}
 (A, X) \mapsto \text{hom}(A, X) & A \xrightarrow{g} X & g \\
 \uparrow u \quad \downarrow f & \downarrow & \downarrow \\
 (B, Y) & \text{hom}(B, Y) & \left(\begin{array}{ccc} B & & Y \\ u \downarrow & & \uparrow f \\ A & \xrightarrow{g} & X \end{array} \right) & f \circ g \circ u
 \end{array}$$

\mathcal{A} , $\text{co}\mathcal{A}$ arise from an op'n on categories which is called

CURRYING

$$\begin{aligned}
 [A \times B, \mathcal{C}] &\cong [A, [B, \mathcal{C}]] \\
 &[B, [A, \mathcal{C}]]
 \end{aligned}$$

To every $F: \mathcal{C}^0 \rightarrow \text{Set}$ associates the category of elements

$$\text{Elt}_s(F) \quad (A, a \in FA) \xrightarrow{h} (B, b \in FB)$$

is $h: A \rightarrow B$ such that $Fh(b) = a$.

F is naturally isom. to $\text{hom}(-, X_F)$ \iff

$\text{Elt}_s(F)$ has a terminal object

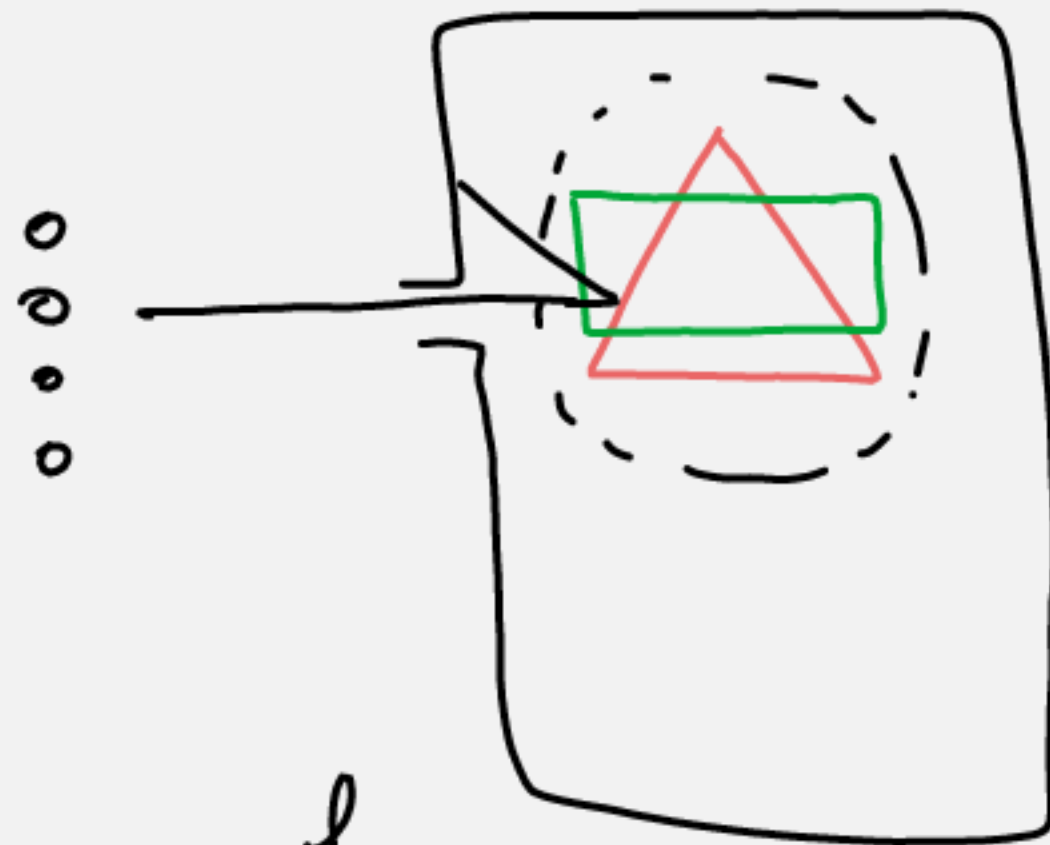
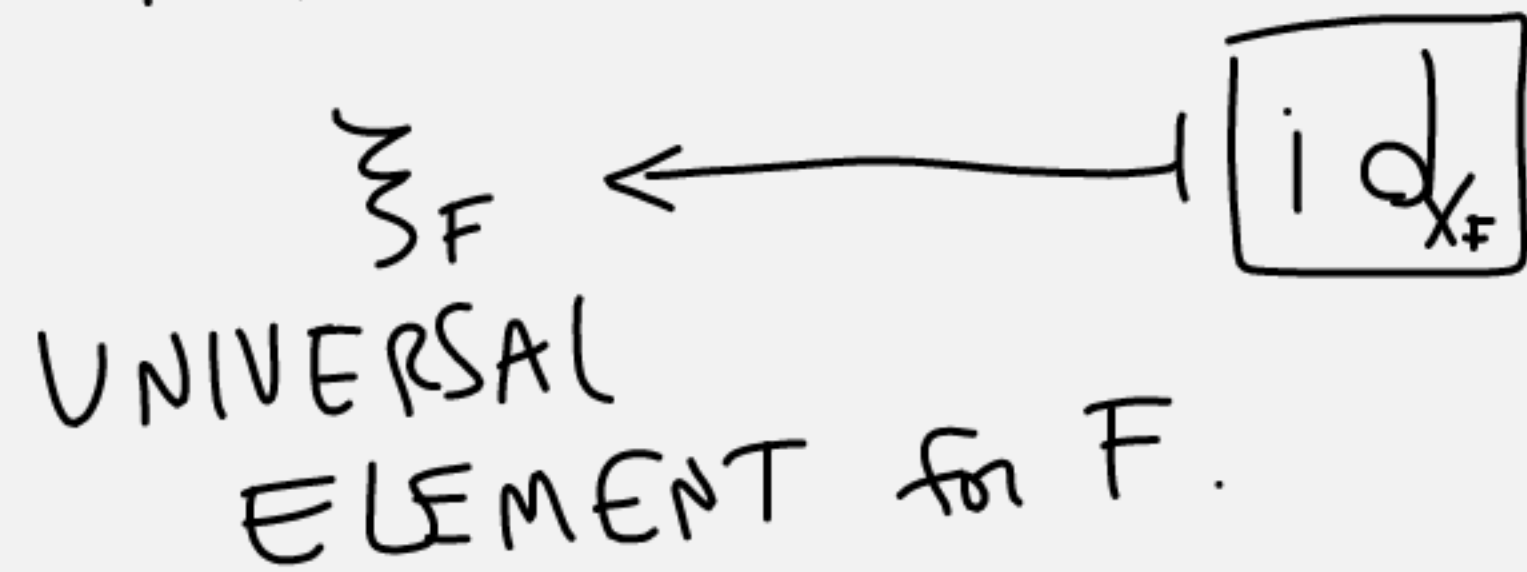
$$FA \cong \text{hom}(A, X_F) \quad (A, a) \xrightarrow{h_{(A,a)!}} (X_F, \sum_{F \in F} X_F)$$

$a \xleftrightarrow{\quad} h_{(A,a)}$

$$\Sigma_F \in F X_F$$

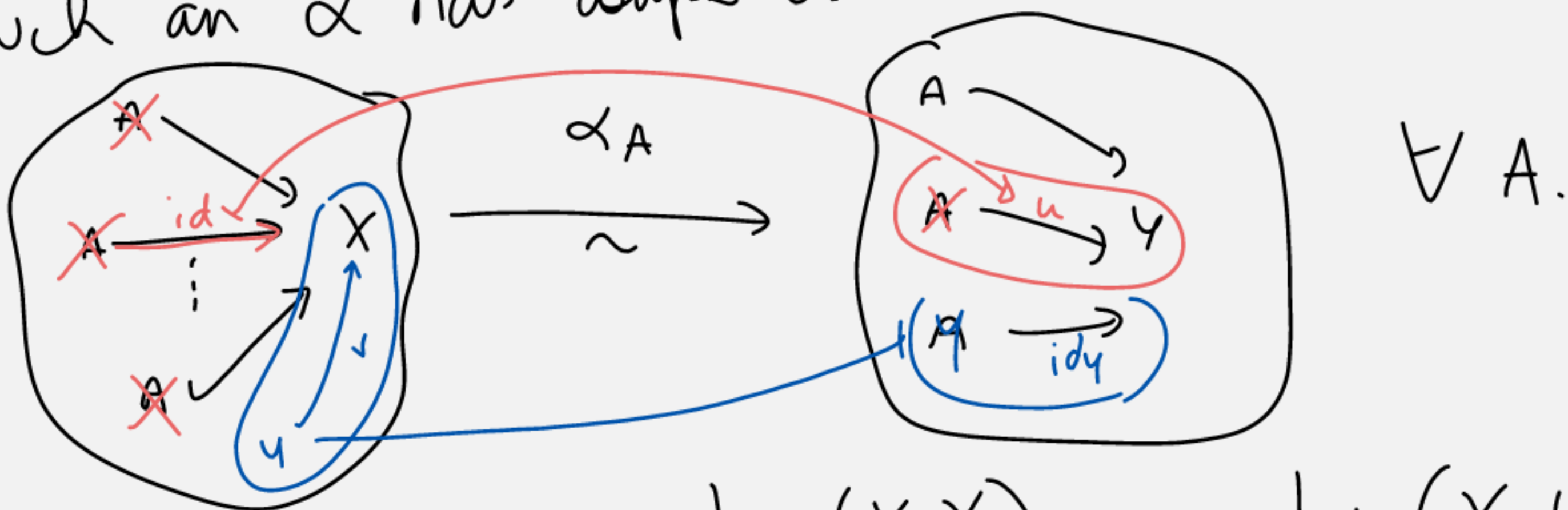
$$F(X_F) \cong \text{hom}(X_F, X_F)$$

arrows $X_F \rightarrow X_F$



To decide whether $X, Y \in \mathcal{C}_0$ are isomorphic, it is enough (necessary & sufficient) to find an isomorphism between the representable functors $\text{hom}(-, X) \xrightarrow[\sim]{\alpha} \text{hom}(-, Y)$

Such an α has components



If $A = X$ $\alpha_X: \text{hom}(X, X) \longrightarrow \text{hom}(X, Y)$
 $\text{id}_X \longmapsto u = \alpha_X(\text{id}_X)$

Similarly $\alpha_A^{-1}: \text{hom}(A, Y) \longrightarrow \text{hom}(A, X)$

$A = Y$ $\alpha_Y^{-1}: \text{hom}(Y, Y) \longrightarrow \text{hom}(Y, X)$
 $\text{id}_Y \longmapsto v = \alpha_Y^{-1}(\text{id}_Y)$

YL I) Natural tns of type $\alpha: \text{hom}(-, X) \Rightarrow F$
 are uniquely determined by the image of id_X
 $\alpha_A: \text{hom}(A, X) \rightarrow FA$ acts in a way that is
 completely determined by $\alpha_X(\text{id}_X) \in FX$
 (This means that given $\alpha, \beta: \text{hom}(-, X) \Rightarrow F$
 $\alpha = \beta$ ($\alpha_A = \beta_A \forall A$)
 IFF $\alpha_X(\text{id}_X) = \beta_X(\text{id}_X$)

YL II) Every element $\xi \in FX$ extends
 to a natural transformation
 $\alpha: \text{hom}(-, X) \Rightarrow F$
 such that $\alpha_X(\text{id}_X) = \xi \in FX$

YL I, YL II prove that
 $\text{Nat}(\text{hom}(-, X), F) \xrightarrow{\mathbf{I}} FX$ is a bijection (YL I injective)
 $\{\alpha_A: \text{hom}(A, X) \rightarrow FA\} \mapsto \alpha_X(\text{id}_X)$ (YL II surjective)

Yoneda Lemma

$\left\{ \begin{array}{l} \text{natural tns} \\ \text{hom}(-, X) \Rightarrow F \end{array} \right\}$ in bijection with FX

In particular (as a corollary)

$F = \text{hom}(-, Y)$ $\text{Nat}(\text{hom}(-, X), \text{hom}(-, Y)) \cong \text{hom}(-, Y)(X)$
 $\text{hom}(X, Y)$

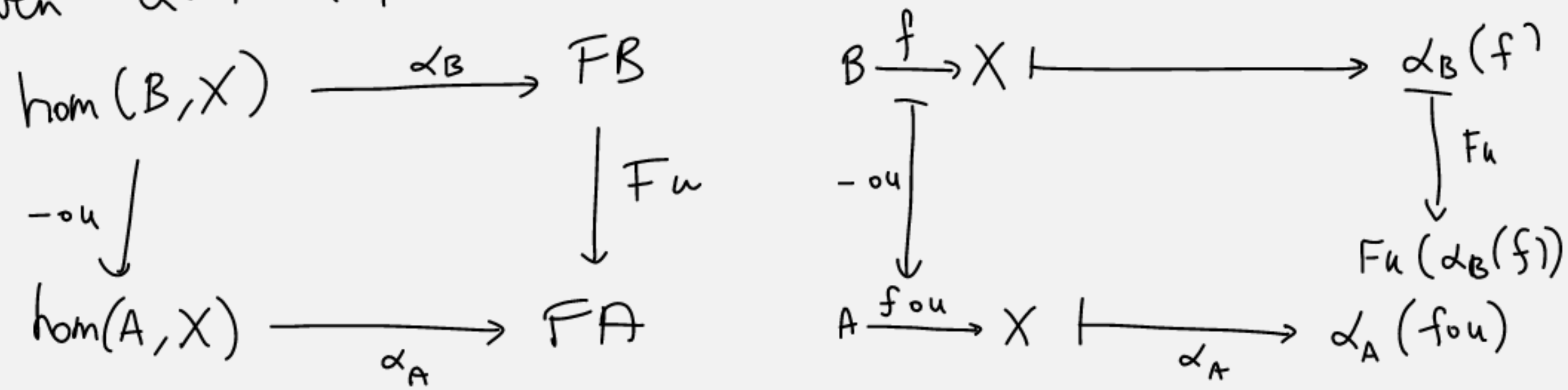
$\text{Nat}(\text{hom}(-, X), \text{hom}(-, Y)) = \text{hom}_{[e^{\circ}, \text{set}]}(\mathcal{L}X, \mathcal{L}Y)$

$\text{hom}_e(X, Y) \cong \text{hom}_{[e^{\circ}, \text{set}]}(\mathcal{L}X, \mathcal{L}Y)$

which is the def'n of fully faithful functor

YL I

Given $\alpha: \text{hom}(-, X) \Rightarrow F$, for every $u: A \rightarrow B$



$\alpha_A(f \circ u) = Fu(\alpha_B(f)) \implies \alpha_A(\text{id}_A \circ u) = Fu(\alpha_X(\text{id}_X))$
 $\forall u: A \rightarrow B$
 $\forall f: B \rightarrow X \quad f = \text{id}_X: X \rightarrow X$

$\alpha_A(\text{id}_A \circ u) = Fu(\alpha_X(\text{id}_X))$
 $\forall u \quad \alpha_A(u) = Fu(\xi_F)$
 $\alpha_A(-) = \text{evaluation at } \xi_F \text{ of } Fu$

The expression $\alpha_A(u) = Fu(\xi_F^\alpha) \quad \alpha_X(\text{id}_X)$

$\beta_A(u)$ (for β of same type $\text{hom}(-, X) \Rightarrow F$)

$\beta_A(u) \equiv Fu(\xi_F^\beta)$
 $\beta_X(\text{id}_X)$

YL II)

Given $\xi \in FX$, have to define a n. transformation

$$\alpha : \text{hom}(-, X) \Rightarrow F$$

$$\alpha_A : \text{hom}(A, X) \longrightarrow FA$$

$$(f: A \rightarrow X) \longmapsto$$

Cook an element of FA from
 $\emptyset - X \in \mathcal{C}_0$

1 - $F: \mathcal{C}^{\text{op}} \rightarrow \text{Set}$,

2 - $\xi: FX$

3 - $f: A \rightarrow X$

$$\alpha_A^{(\xi)}(f) = F(f)(\xi)$$

*btw
by definition*

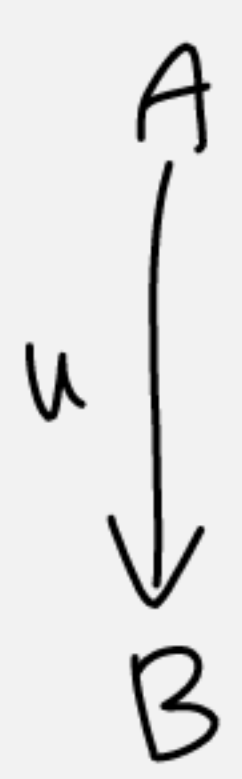
$$\alpha_X^{(\xi)}(\text{id}_X) = F(\text{id}_X)(\xi)$$
$$= \text{id}_{F(X)}(\xi)$$
$$= \xi$$

$$\underline{\alpha_A^{(\xi)}} = \underline{\lambda f. F(f)(\xi)}$$

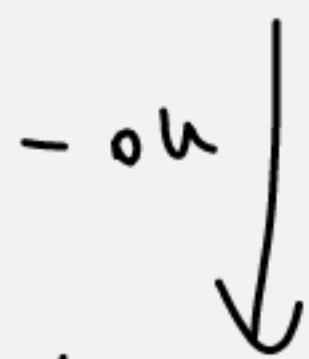
$$FX \xrightarrow{F(f)} FA$$

$$\xi \longmapsto F(f)(\xi)$$

$\{\alpha_A\}$ is a nat. transformation:



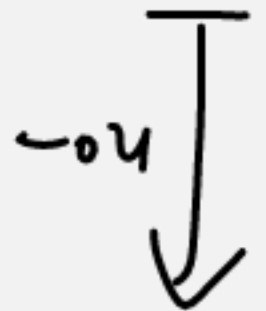
$$\text{hom}(B, X) \xrightarrow{\alpha_B(\xi)} FB$$



$$\text{hom}(A, X) \xrightarrow{\alpha_A(\xi)} FA$$



$$(B \xrightarrow{f} X) \longmapsto F(f)(\xi)$$



$$(A \xrightarrow{f \circ u} X) \longmapsto$$

$$Fu(F(f)(\xi))$$

$$F(f \circ u)(\xi)$$

$$= (Fu \circ Ff)(\xi)$$

(P, \leq) partially ordered set

$[P^\circ, \text{Set}]$ $P^\circ \rightarrow \text{Set}$

$P \rightarrow [P^\circ, \text{Set}]$

$x \mapsto \text{hom}_P(-, x)$
 $a \mapsto \text{hom}(a, x)$ $\begin{cases} \{a \leq x\} & \text{if } a \leq x \\ \emptyset & \text{otherwise} \end{cases}$
 "initial segment" of x

YL says $\text{Nat}(\text{hom}(-, x), F) \cong F(x)$

$\alpha_a: \text{hom}(a, x) \rightarrow Fa$
 $\alpha_a: \emptyset \rightarrow Fa$ \swarrow $a \not\leq x$
 \searrow $a \leq x$ $\{*\} \rightarrow Fa$

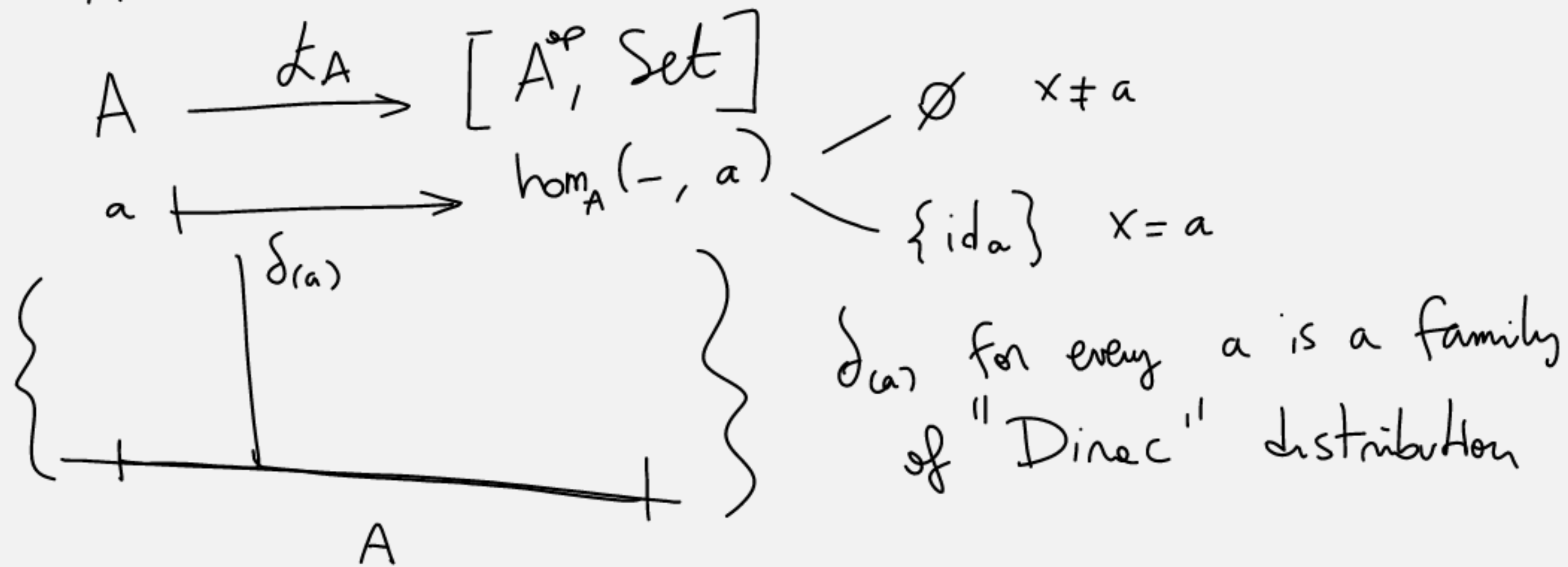
$\alpha: \text{hom}(-, x) \rightarrow \text{hom}(-, y)$
 $\alpha_a: \text{hom}(a, x) \rightarrow \text{hom}(a, y)$
 if $a \leq x$ then $a \leq y$

$\alpha_x(x \leq x) \in \text{hom}(x, y)$ saying that $x \leq y$
 $\alpha_y(y \leq y) \in \text{hom}(y, x)$ " " $y \leq x$

$\alpha^{-1}: \text{hom}(-, y) \rightarrow \text{hom}(-, x)$
 is the converse
 if $a \leq y$ then $a \leq x$

$\Rightarrow x = y$

A set (discrete category)
 $A^{\text{op}} \rightarrow \text{Set}$ amounts to a family of sets $\{F_a \mid a \in A\}$



YL says that a generic "distribution" $A^{\text{op}} \xrightarrow{F} \text{Set}$

is reconstructed from

$$\{ \delta(a) \implies F \} = F_a$$

$$\{ \delta(a)(x) \implies F(x) \}$$

$$\{ \emptyset \text{ } a \neq x, \{ \delta_x \} \implies F(a) \} \equiv F(a)$$

